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# Subword complexes via triangulations of root polytopes

### Laura Escobar & Karola Mészáros

ABSTRACT Subword complexes are simplicial complexes introduced by Knutson and Miller to illustrate the combinatorics of Schubert polynomials and determinantal ideals. They proved that any subword complex is homeomorphic to a ball or a sphere and asked about their geometric realizations. We show that a family of subword complexes can be realized geometrically via regular triangulations of root polytopes. This implies that a family of  $\beta$ -Grothendieck polynomials are special cases of reduced forms in the subdivision algebra of root polytopes. We can also write the volume and Ehrhart series of root polytopes in terms of  $\beta$ -Grothendieck polynomials.

### 1. INTRODUCTION

In this paper we provide geometric realizations of pipe dream complexes  $PD(\pi)$  of permutations  $\pi = 1\pi'$ , where  $\pi'$  is a dominant permutation on 2, 3, ..., n as well as the subword complexes that are the cores of the pipe dream complexes  $PD(\pi)$ . We realize  $PD(\pi)$  as (repeated cones of) regular triangulations of the root polytopes  $\mathcal{P}(T(\pi))$ .

Subword complexes are simplicial complexes introduced by Knutson and Miller in [10, 11] to illustrate the combinatorics of Schubert polynomials and determinantal ideals. Since the appearance of Knutson's and Miller's work in there has been a flurry of research into the geometric realization of subword complexes with progress in realizing families of spherical subword complexes: [2, 3, 4, 5, 20, 21, 22, 24]. This paper is the first to succeed in realizing a family of subword complexes which are homeomorphic to balls.

Subword complexes were first shown to relate to triangulations of root polytopes by Mészáros in [17], where the author gives a geometric realization of the pipe dream complex of [1, n, n - 1, ..., 2] and whose work served as the stepping stone for the present project. In the papers [12, 14, 15, 16] Mészáros studied triangulations of root polytopes that we utilize in this work (some of the mentioned papers are in the language of flow polytopes, but in view of [17, Section 4] some of their content can also be understood in the language of root polytopes).

The main theorem of this paper is the following, which has several interesting consequences explored in the paper. For the definitions needed for this theorem see the later sections.

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THEOREM 1.1. Let  $\pi = 1\pi' \in S_n$ , where  $\pi'$  is dominant. Let  $C^2(\pi)$  be the core of  $PD(\pi)$  coned over twice. The canonical triangulation of the root polytope  $\mathcal{P}(T(\pi))$  (which is a regular triangulation) is a geometric realization of  $C^2(\pi)$ .

The roadmap of this paper is as follows. In Sections 2 and 3 we explain the necessary background about subword complexes and root polytopes, respectively. In Section 4 we prove a geometric realizations of pipe dream complexes  $PD(\pi)$  of permutations  $\pi = 1\pi'$ , where  $\pi'$  is a dominant permutation on  $2, 3, \ldots, n$ , via triangulations of root polytopes  $\mathcal{P}(T(\pi))$ . In Section 5 we use the previous result to show that  $\beta$ -Grothendieck polynomials are special cases of reduced forms in the subdivision algebra of root polytopes while in Section 6 we show how to express the volume and Ehrhart series of root polytopes in terms of Grothendieck polynomials. Appendix A is devoted to proving a certain uniqueness property of the reduced form in the subdivision algebra that we used in Section 5.

### 2. Background on pipe dream complexes

We let  $S_n$  denote the set of permutations of size n.

DEFINITION 2.1. The (Rothe) diagram of a permutation  $\pi \in S_n$  is the collection of boxes  $D(\pi) = \{(\pi_j, i) : i < j, \pi_i > \pi_j\}$ . It can be visualized by considering the boxes left in the  $n \times n$  grid after we cross out the boxes appearing south and east of each 1 in the permutation matrix for  $\pi$ .



FIGURE 1. The diagram for  $\pi = [4132]$ .

Notice that no two permutations can give the same diagram. We will consider permutations of the form  $\pi = 1\pi'$  where  $\pi'$  is a *dominant permutation* of  $\{2, \ldots, n\}$ , i.e., the diagram of  $\pi$  is a partition with north-west most box at position (2, 2). Dominant permutations can be equivalently defined as the 132-avoiding permutations, and there are Catalan many for fixed size. Our convention is to encode the partition by the number of boxes in each column.



FIGURE 2. The diagram for  $\pi = [164235]$  which corresponds to  $\lambda = (4, 2)$ 

DEFINITION 2.2. A pipe dream for  $\pi \in S_n$  is a tiling of an  $n \times n$  matrix with two tiles, crosses + and elbows  $\neg$ , such that

(1) all tiles in the weak south-east triangle of the  $n \times n$  matrix are elbows, and

(2) if we write 1, 2, ..., n on the left and follow the strands (ignoring second crossings among the same strands) they come out on the top and read π.

A pipe dream is reduced if no two strands cross twice.



FIGURE 3. A reduced pipe dream for  $\pi = [1423]$ .

DEFINITION 2.3. The pipe dream complex  $PD(\pi)$  of a permutation  $\pi \in S_n$  is the simplicial complex with vertices given by entries on the northwest triangle of an  $n \times n$ -matrix and facets given by the elbow positions in the reduced pipe dreams for  $\pi$ .

Pipe dream complexes are a special case of the subword complexes defined by Knutson and Miller in [10, 11]. We proceed to explain the correspondence. The group  $S_n$  is generated by the adjacent transpositions  $s_1, \ldots, s_{n-1}$ , where  $s_i$  transposes  $i \leftrightarrow i+1$ . Let  $Q = (q_1, \ldots, q_m)$  be a word in  $\{s_1, \ldots, s_{n-1}\}$ , i.e., Q is an ordered sequence. A subword  $J = (r_1, \ldots, r_m)$  of Q is a word obtained from Q by replacing some of its letters by -. There are a total of  $2^{|Q|}$  subwords of Q. Given a subword J, we denote by  $Q \smallsetminus J$  the subword with k-th entry equal to - if  $r_k \neq -$  and equal to  $q_k$  otherwise for,  $k = 1, \ldots, m$ . For example,  $J = (s_1, -, s_3, -, s_2)$  is a subword of  $Q = (s_1, s_2, s_3, s_1, s_2)$ and  $Q \searrow J = (-, s_2, -, s_1, -)$ . Given a subword J we denote by  $\prod J$  the product of the letters in J, from left to right, with - behaving as the identity.

DEFINITION 2.4 ([10, 11]). Let  $Q = (q_1, \ldots, q_m)$  be a word in  $\{s_1, \ldots, s_{n-1}\}$  and  $\pi \in S_n$ . The subword complex  $\Delta(Q, \pi)$  is the simplicial complex on the vertex set Q whose facets are the subwords F of Q such that the product  $\prod(Q \setminus F)$  is a reduced expression for  $\pi$ .

In this language,  $PD(\pi)$  is the subword complex  $\Delta(Q, \pi)$  corresponding to the triangular word  $Q = (s_{n-1}, s_{n-2}, s_{n-1}, \ldots, s_1, s_2, \ldots, s_{n-1})$  and  $\pi$ . The correspondence between pipe dreams and subwords is induced by the labeling of the entries in the northwest triangle of an  $n \times n$ -matrix by adjacent transpositions, as depicted in Figure 4, and by making a + in a pipe dream correspond to a - in a subword and a - correspond to the  $s_i$  in its entry. In order to go from a pipe dream to a subword, we read the entries in the northwest triangle from left to right starting at the bottom.



FIGURE 4. Labeling of the entries in the northwest triangle by adjacent transpositions.

DEFINITION 2.5. Let  $cone(\pi)$  be the set of vertices of  $PD(\pi)$  that are in all its facets. We define the core of  $\pi$ , denoted by  $core(\pi)$ , to be the simplicial complex obtained by restricting  $PD(\pi)$  to the set of vertices not in  $cone(\pi)$ .

Notice that  $PD(\pi)$  is obtained from its core by iteratively coning the simplicial complex  $\operatorname{core}(\pi)$  over the vertices in  $\operatorname{cone}(\pi)$ . This is a standard definition for simplicial complexes. In the language of pipe dream complexes, the core is the restriction to the entries in the  $n \times n$  matrix that are a cross in some reduced pipe dream for  $\pi$ . Following the correspondence described in Figure 4, this restriction induces a subword Q' of the triangular word and so  $\operatorname{core}(\pi)$  is the subword complex  $\Delta(Q', \pi)$ .

EXAMPLE 2.6. Figure 3 shows a reduced pipe dream for  $\pi = [1423]$ . The pipe dream complex PD(1423) is three dimensional with vertices (1, 1), (1, 2), (1, 3), (2, 1), (2, 2), and (3, 1). Every facet contains the vertices (1, 1) and (1, 3) since every reduced pipe dream must have elbows on these positions. Therefore, one can recover PD(1423) from its core, shown in Figure 5, by coning twice.



FIGURE 5. The core of PD(1423). The facets are labelled by the reduced pipe dreams for [1423].

Since we are only considering permutations of the form  $1\pi'$  with  $\pi'$  a dominant permutation,  $\operatorname{core}(1\pi')$  is easy to describe. Given a diagram of a permutation there are two natural reduced pipe dreams for  $\pi$ , referred to as the *bottom reduced pipe dream of*  $\pi$  and the *top reduced pipe dream of*  $\pi$ , one obtained by aligning the diagram to the left and replacing the boxes with crosses and the other one by aligning the diagram up. See Figure 6.



(A) Aligned left is the botton reduced pipe dream

(B) Aligned up is the top reduced pipe dream  $\ensuremath{\mathsf{B}}$ 

FIGURE 6. Two reduced pipe dreams for [164235] obtained by aligning the diagram to the left and to the top.

The core of  $1\pi'$  is the simplicial complex obtained by restricting  $PD(\pi)$  to the vertices corresponding to the positions of the crosses in the superimposition of these

two pipe dreams. We refer to the region itself as the *core region*, and denote it by  $cr(\pi)$ . See Figure 7 for an example. Note that different permutations can have the same core region, as is the case for [15342] and [15432].



FIGURE 7. The core region of [164235]

In [1], Bergeron and Billey introduced an algorithm to construct all reduced pipe dreams for  $\pi$ . Given a reduced pipe dream P for all permutations  $\pi$ , a ladder admitting rectangle is a connected  $k \times 2$  rectangle inside P such that  $k \ge 2$  and the only  $\neg_{\frown}$  inside this rectangle are in the top row and in the southeast corner, see the diagram on the left in Figure 8. A *ladder move* on P moves the  $\neg_{\frown}$  in the southwest corner of a ladder admitting rectangle to the northeast corner. Notice that the resulting pipe dream is a reduced pipe dream for  $\pi$ .



FIGURE 8. Ladder move.

THEOREM 2.7 ([1]). The set of all reduced pipe dreams of  $\pi$  equals the set of pipe dreams that can be derived from the bottom reduced pipe dream by a sequence of ladder moves.

The boundary of a pure simplicial complex  $\Delta$  is the simplicial complex  $\partial \Delta$  with facets the codimension 1 faces of  $\Delta$  that are in exactly one facet of  $\Delta$ . A face F of  $\Delta$  is *interior* if F is not in  $\partial \Delta$ .

### 3. BACKGROUND ON ROOT POLYTOPES

We follow the exposition of [17, Section 4] in this section. A root polytope of type  $A_n$  is the convex hull of the origin and some points in  $\Phi^+ = \{e_i - e_j \mid 1 \leq i < j \leq n+1\}$ , the set positive roots of type  $A_n$ , where  $e_i$  denotes the  $i^{th}$  coordinate vector in  $\mathbb{R}^{n+1}$ . An important family of root polytopes studied by Gelfand, Graev and Postnikov in [7] are the full root polytopes

$$\mathcal{P}(A_n^+) = \text{ConvHull}(0, e_i - e_j \mid 1 \leq i < j \leq n+1).$$

In this paper we restrict ourselves to a class of root polytopes including  $\mathcal{P}(A_n^+)$ , which have subdivision algebras as defined in [12]. We discuss subdivision algebras in relation to Grothendieck polynomials in Section 5.

Let G be a forest on the vertex set [n+1]. Define

- V<sub>G</sub> = {e<sub>i</sub> e<sub>j</sub> | (i, j) ∈ E(G), i < j}, a set of vectors associated to G;</li>
  cone(G) = ⟨V<sub>G</sub>⟩ := {∑<sub>e<sub>i</sub>-e<sub>j</sub>∈V<sub>G</sub> c<sub>ij</sub>(e<sub>i</sub> e<sub>j</sub>) | c<sub>ij</sub> ≥ 0}, the cone associated to
  </sub> G; and
- $\overline{\mathcal{V}}_G = \Phi^+ \cap \operatorname{cone}(G)$ , all the positive roots of type  $A_n$  contained in  $\operatorname{cone}(G)$ . The root polytope  $\mathcal{P}(G)$  associated to the forest G is

(3.1) 
$$\mathcal{P}(G) := \text{ConvHull}(0, e_i - e_j \mid e_i - e_j \in \overline{\mathcal{V}}_G).$$

The root polytope  $\mathcal{P}(G)$  associated to a graph G can also be defined as

(3.2) 
$$\mathcal{P}(G) = \mathcal{P}(A_n^+) \cap \operatorname{cone}(G).$$

Note that  $\mathcal{P}(A_n^+) = \mathcal{P}(P_{n+1})$  for the path graph  $P_{n+1}$  on the vertex set [n+1].

The reduction rule for graphs: Given a graph  $G_0$  on the vertex set [n + 1] and  $(i,j), (j,k) \in E(G_0)$  for some i < j < k, let  $G_1, G_2, G_3$  be graphs on the vertex set [n+1] with edge sets

(3.3)  

$$E(G_1) = E(G_0) \setminus \{(j,k)\} \cup \{(i,k)\},$$

$$E(G_2) = E(G_0) \setminus \{(i,j)\} \cup \{(i,k)\},$$

$$E(G_3) = E(G_0) \setminus \{(i,j), (j,k)\} \cup \{(i,k)\}$$

We say that  $G_0$  reduces to  $G_1, G_2, G_3$  under the reduction rules defined by equations (3.3).

LEMMA 3.1 (Reduction Lemma for Root Polytopes, [12]). Given a forest  $G_0$  with d edges let  $(i, j), (j, k) \in E(G_0)$  for some i < j < k and  $G_1, G_2, G_3$  as described by equations (3.3). Then

$$\mathcal{P}(G_0) = \mathcal{P}(G_1) \cup \mathcal{P}(G_2)$$
  
where all polytopes  $\mathcal{P}(G_0), \mathcal{P}(G_1), \mathcal{P}(G_2)$  are d-dimensional and  
 $\mathcal{P}(G_3) = \mathcal{P}(G_1) \cap \mathcal{P}(G_2)$  is  $(d-1)$ -dimensional.

The Reduction Lemma says that performing a reduction on a forest  $G_0$  is the same as dissecting the d-dimensional polytope  $\mathcal{P}(G_0)$  into two d-dimensional polytopes  $\mathcal{P}(G_1)$  and  $\mathcal{P}(G_2)$ , whose vertex sets are subsets of the vertex set of  $\mathcal{P}(G_0)$ , whose interiors are disjoint, whose union is  $\mathcal{P}(G_0)$ , and whose intersection is a facet of both.

The following theorem in [12] describes a triangulation of the root polytope  $\mathcal{P}(G)$ for any forest G. Its proof is based on the Reduction Lemma stated above. We now define the terminology used in the theorem. A graph G on the vertex set [n+1] is said to be *noncrossing* if there are no  $1 \le i < j < k < l \le n+1$  such that (i,k), (j,l)are edges of G. The graph G is said to be *alternating* if at each vertex v of G all edges are either of the form (v,i) for v < i or of the form (i,v) for i < v. Finally, the directed transitive closure of the graph G on the vertex set [n + 1], denoted by  $\overline{G}$ , is  $\overline{G} := (V(G), E(G) \cup \{(i, j) \mid (i, j) \notin E(G) \text{ and there exist } i < i_1 < \cdots < i_k < i_$ *j* with  $(i, i_1), (i_1, i_2), \dots, (i_k, j) \in E(G)$ }.

THEOREM 3.2 ([12]). Let  $T_1, \ldots, T_k$  be the noncrossing alternating spanning trees of the directed transitive closure of the noncrossing forest G. Then  $\mathcal{P}(T_1), \ldots, \mathcal{P}(T_k)$  are top dimensional simplices in a regular triangulation of  $\mathcal{P}(G)$ .

We note that there is a version of Theorem 3.2 in [12] that does not require the noncrossing condition on G; however, in the present paper we only invoke the above version which has the advantage that it is easier to state.

We refer to the triangulation specified in Theorem 3.2 as the *canonical triangulation* of  $\mathcal{P}(G)$ . We remark that the polytopes  $\mathcal{P}(T_i)$ ,  $i \in [k]$ , in Theorem 3.2 are simplices, because the graphs  $T_i$ ,  $i \in [k]$ , are alternating trees, and as such they are their own

transitive closure, with the roots corresponding to the edges of each of them linearly independent. Since each simplex  $\mathcal{P}(T_i)$ ,  $i \in [k]$ , contains the vertex 0, it follows that the canonical triangulation of  $\mathcal{P}(G)$  also induces a triangulation of the vertex figure of  $\mathcal{P}(G)$  at 0, which we also call the canonical triangulation of the vertex figure of  $\mathcal{P}(G)$  at 0. We sumarize some facts about the canonical triangulation in the following proposition.

PROPOSITION 3.3. Let  $\mathcal{C}(G)$  denote the simplicial complex induced by the canonical triangulation of the vertex figure of  $\mathcal{P}(G)$  at 0. The vertices of  $\mathcal{C}(G)$  are in bijection with edges (i, j) in the directed transitive closure of G; the vertex of  $\mathcal{C}(G)$  corresponding to (i, j) is the intersection of the ray pointing to  $e_i - e_j$  and the hyperplane by which we intersect  $\mathcal{P}(G)$  to obtain the considered vertex figure.

4. PIPE DREAM COMPLEXES AS TRIANGULATIONS OF ROOT POLYTOPES

In this section we give geometric realizations of pipe dream complexes of permutations  $\pi = 1\pi'$ , where  $\pi'$  is dominant, in terms of triangulations of root polytopes. Indeed, we construct a geometric realization of the subword complex that is the core of  $PD(1\pi')$ . To this end we start by defining a tree  $T(\pi)$  for each permutation  $\pi = 1\pi'$ ,  $\pi'$  dominant.

4.1. CONSTRUCTION OF THE TREE  $T(\pi)$ . Let  $\pi = 1\pi'$ , where  $\pi'$  is dominant. Denote by  $S(\pi)$  the subword complex that is the core $(\pi)$  coned over the vertex of  $PD(\pi)$ corresponding to the entry (1, 1). Denote the region that is the union of (1, 1) and  $cr(\pi)$ by  $R(\pi)$ . In order to determine the tree  $T(\pi)$ , we will label the southeast boundary with some numbers and we will place dots in some entries of  $R(\pi)$ , see Figure 9. The boundary of the core region starting from the southwest (SW) corner of it to the northeast (NE) corner can be described as a series of east (E) and north (N) steps. Let A be the set consisting of all the N steps together with some E steps. The step  $E \in A$ if the bottom reduced pipe dream is bounded by E but not by the N step directly preceding E. As we traverse this lower boundary from the SW corner we write the numbers  $1, \ldots, m$  in increasing fashion below the E steps and to the right of the N steps that belong to A. For the E steps that we did not assign a number, we consider their number to be the number assigned to the N step directly preceding them.

We now describe how to place dots in  $R(\pi)$ . Consider the bottom reduced pipe dream drawn inside  $R(\pi)$  and with elbows replaced by dots. Drop these dots south. Define  $T(\pi)$  to be the tree on m vertices such that there is an edge between vertices i < j if there is a dot in the entry in the column of the E step labeled i and in the row of the N step labeled j. Let  $t(\pi)$  be the number of edges of  $T(\pi)$ .



FIGURE 9. Let  $\pi = [15342]$ , on the left we see the bottom reduced pipe dream for  $\pi$  drawn inside  $R(\pi)$  with dots instead of elbows, this gives the labeling of the boundary. We then drop the dots to the south to get the dots encoding  $T(\pi)$ , which is depicted on the right.

LEMMA 4.1. For permutations  $\pi = 1\pi'$ , where  $\pi'$  is dominant, the graph  $T(\pi)$  constructed above is a tree.

*Proof.* Our proof that  $T(\pi)$  is acylic and connected is based on the following observations which we number for convenience.

Observation 1. Consider the bottom reduced pipe dream drawn inside  $R(\pi)$  and with elbows replaced by dots as in the leftmost diagram in Figure 9. Each box in the top row of  $R(\pi)$  contains a dot. In every other row of  $R(\pi)$  there is either zero or one dot depending on whether or not the N step on the right of the considered row is adjacent to the bottom pipe dream or not.

Observation 2. The only dots that can change position when dropping them in the above construction are the ones on the top row of  $R(\pi)$ .

First we show that  $T(\pi)$  is acylic. Note that after dropping the dots south, if there is a dot in the box (i, j) of  $R(\pi)$  then there are dots on all positions (k, l) with  $k \ge i$ or  $j \ge l$ . Indeed, if  $T(\pi)$  has a cycle then the labelled southeast boundary of  $R(\pi)$ together with the dropped dots must have a piece that looks like:



Consider the step before dropping the dots in  $R(\pi)$ . By Observation 2, there is a crossing on the box on top of the box bounded by the N step for b. In order for b to be the label of both the N and E step noted in the above figure, we conclude that before dropping the dots in  $R(\pi)$  we must have the following confirguration:



However, this contradicts the very construction of  $R(\pi)$ . Therefore,  $T(\pi)$  is acyclic.

To show that  $T(\pi)$  is a tree, since it is acyclic it suffices to show that it has  $t(\pi) = m - 1$  edges, where m is the number of distinct labels used in the boundary.

By construction, the number of dots in  $R(\pi)$  is  $m_1 + m_2 - 1$  where  $m_1 = l(\lambda) + 1$  is the number of dots in the first row before they are dropped, and  $m_2$  is the number of N steps in the boundary not adjacent to the bottom pipe dream. Observation 1 and the fact that the last dot in the first row is to the left of an N step yields that the number of dots in  $R(\pi)$ , which is also  $t(\pi)$ , is  $m_1 + m_2 - 1$ .

On the other hand, note that m, the number of labels equals  $m_1 + m_2$  concluding our proof. Indeed, the number of E steps along the boundary equals  $m_1$  and they all have distinct labels. Moreover, there are exactly  $m_2$  N steps that have labels not occurring among the labels of the E steps.

REMARK 4.2. We note that the construction of  $T(\pi)$  can be simplified in the special case of  $\pi = 1\pi'$  with  $D(\pi')$  a partition with distinct parts. Indeed, in the above construction the first string of E steps are in A. Furthermore, the E steps that can be seen as the boundary of the bottom reduced pipe dream B of  $\pi$  and such that E

bounds row r of length  $l_r$  of B and the row below row r of length  $l_{r+1}$  is at least two boxes shorter than row r and moreover, step E is not  $(l_{r+1} + 1)^{st}$  from the left side are also in A. Also, we are placing dots in the rightmost boxes of the core region as well as in positions  $(l_{r+1} + 1)^{st}$  until  $(l_r - 1)^{th}$  in rows r of the bottom reduced pipe dream B that are longer than row r + 1 by at least two boxes. In the case in which  $\pi = 1\pi'$  where  $\pi'$  is dominant with its diagram having all parts distinct, then the decoration on the core diagram is much simpler. The first string of E steps consist of only one E step and this is also the only E step in A, so the number of vertices of  $T(\pi)$  is 2 more than the size of the largest column of the diagram of  $\pi$ . The dots are placed on the rightmost boxes of  $cr(\pi)$ , see Figure 10.



FIGURE 10. Obtaining  $T(\pi)$  from  $R(\pi)$ , where  $\pi = [164235]$ .

4.2. BIJECTION BETWEEN THE VERTICES OF  $S(\pi)$  AND THE VERTICES OF  $C(T(\pi))$ . The vertices of  $S(\pi)$  are in bijection with configurations of one elbow and  $|R(\pi)| - 1$ crosses in  $R(\pi)$ , where  $|R(\pi)|$  equals the number of entries in  $R(\pi)$ . Denote these vertices by  $v_1, \ldots, v_k$ . We define a map M from the vertices of the simplicial complex  $S(\pi)$  to the vertices of  $C(\pi) := C(T(\pi))$ . Recall that  $C(\pi)$  is the canonical triangulation of the vertex figure at 0 of the root polytope  $\mathcal{P}(T(\pi))$  and by Proposition 3.3 the vertices of the triangulation are in bijection with edges (i, j) in the directed transitive closure of  $T(\pi)$ . The latter in turn are in bijection with the boxes of  $R(\pi)$  by the map that takes a box to the edge (i, j) if the E step below the box and in the boundary of  $R(\pi)$  is labeled by *i* and the N step to the right of the box and in the boundary of  $R(\pi)$  is labeled by *j*. The map M is defined analogously as follows.

DEFINITION 4.3. Consider a vertex of  $S(\pi)$ ; this can be seen as a sole elbow tile in  $R(\pi)$ . M maps this vertex to the vertex of  $C(\pi)$  corresponding to (i, j) if the box containing the elbow tile yields the edge (i, j) in  $T(\pi)$  (that is to the intersection of the ray pointing to  $e_i - e_j$  and the hyperplane by which we intersect  $\mathcal{P}(T(\pi))$  to obtain the considered vertex figure).

THEOREM 4.4. The triangulation  $C(\pi)$  is a geometric realization of the subword complex  $S(\pi)$ .

*Proof.* We show that the map M described above respects the simplicial complex structure of  $S(\pi)$  and  $C(\pi)$ . Since both  $S(\pi)$  and  $C(\pi)$  are pure simplicial complexes of the same dimension by Lemma 4.7 it suffices to show that the map M is a bijection on the facets of  $S(\pi)$  and  $C(\pi)$ . This is proven in Theorem 4.8.

Proof of Theorem 1.1. Since  $\mathcal{C}(\pi)$  is the canonical triangulation of the vertex figure at 0 of the root polytope  $\mathcal{P}(T(\pi))$  then we must cone this triangulation once to obtain the canonical triangulation of  $\mathcal{P}(T(\pi))$ . By Theorem 4.4 it follows that this triangulation of  $\mathcal{P}(T(\pi))$  is a geometric realization of the subword complex  $\mathcal{S}(\pi)$ , coned over once. This simplicial complex equals  $\mathcal{C}^2(\pi)$ .



FIGURE 11. A few examples of the image of M. Each colored dot represents the vertex with an elbow at that position.

4.3. GEOMETRIC REALIZATION OF  $CORE(\pi)$ . Next we show that from Theorem 4.4 it follows that we can also realize  $core(\pi)$  geometrically, which is a subword complex as explained in Section 2 after Definition 2.5. To this end we prove an auxiliary lemma first.

LEMMA 4.5. Let  $\pi = 1\pi'$ , with  $\pi'$  dominant. If  $\mathcal{C}(\pi)$  has an interior vertex, then it is the unique vertex in  $\mathcal{C}(\pi)$  on the ray between 0 and  $e_1 - e_m$ . Moreover,  $\mathcal{C}(\pi)$  has an interior vertex if and only if  $\pi = 1n(n-1)\dots 2$ .

Proof.  $\mathcal{C}(\pi)$  has an interior vertex if and only if the cone generated by  $e_i - e_j$  for  $(i,j) \in T(\pi)$  has an interior point  $e_k - e_l$ , where (k,l) is in the directed transitive closure of  $T(\pi)$ . Since  $e_i - e_j$  for  $(i,j) \in T(\pi)$  are linearly independent, an interior point can be expressed as  $\sum_{(i,j)\in T(\pi)} c_{ij}(e_i - e_j)$  with  $c_{ij} > 0$ ,  $(i,j) \in T(\pi)$ . If  $T(\pi) = ([m], \{(i,i+1) \mid i \in [m-1]\})$ , then  $e_1 - e_m = \sum_{(i,j)\in T(\pi)} (e_i - e_j)$  is an interior point; moreover,  $e_k - e_l$  for  $1 \leq k < l \leq m$  is an interior point if and only if k = 1 and l = m. We have  $T(\pi) = ([m], \{(i,i+1) \mid i \in [m-1]\})$  exactly for  $\pi = 1m(m-1)\dots 2$ . For  $T(\pi) \neq ([m], \{(i,i+1) \mid i \in [m-1]\})$ , there is no (k,l) in the directed transitive closure of  $T(\pi)$  such that  $e_k - e_l = \sum_{(i,j)\in T(\pi)} c_{ij}(e_i - e_j)$  with  $c_{ij} > 0$ .

THEOREM 4.6. Let  $\pi = 1\pi'$ , with  $\pi'$  dominant. Let v be the unique vertex in  $C(\pi)$ on the ray between 0 and  $e_1 - e_m$ . For  $\pi \neq 1m(m-1)\ldots 2$ , v is in the boundary of  $C(\pi)$ , and  $\operatorname{core}(\pi)$  is realized by the induced triangulation of the vertex figure of  $C(\pi)$ at v. For  $\pi = 1m(m-1)\ldots 2$ ,  $\operatorname{core}(\pi)$  is realized by the induced triangulation of the boundary of  $C(\pi)$ .

Proof. The vertex v, which is the unique vertex in  $\mathcal{C}(\pi)$  on the ray between 0 and  $e_1 - e_m$ , is the unique coning point of the geometric realization of  $\mathcal{S}(\pi)$ . If v is in the boundary of  $\mathcal{C}(\pi)$ , which happens exactly when  $\pi \neq 1m(m-1)\dots 2$  by Lemma 4.5, then the induced triangulation of the vertex figure at v of  $\mathcal{C}(\pi)$  is a geometric realization of  $\operatorname{core}(\pi)$  (which is homeomorphic to a ball). For  $\pi = 1m(m-1)\dots 2$ , the coning point v lies in the interior of  $\mathcal{C}(\pi)$ , then since it is the only point in the interior of the canonical triangulation  $\mathcal{C}(\pi)$  by Lemma 4.5, then the induced triangulation of the sequence of  $(\pi)$  (which is homeomorphic to a ball). For  $\pi = 1m(m-1)\dots 2$ , the coning point v lies in the interior of  $\mathcal{C}(\pi)$ , then since it is the only point in the interior of the boundary of  $\mathcal{C}(\pi)$  is a geometric realization of  $\operatorname{core}(\pi)$  (which is homeomorphic to a sphere).

We now proceed to prove an auxiliary lemma used in the proof of Theorem 4.4.

LEMMA 4.7. The core of the pipe dream complex of  $\pi = 1\pi'$ , where the diagram of  $\pi'$ is a partition, is of dimension  $t(\pi) - 2$ . The dimension of the root polytope  $\mathcal{P}(T(\pi))$ is  $t(\pi)$  and its vertex figure at 0 is of dimension  $t(\pi) - 1$ . In particular, both  $\mathcal{S}(\pi)$  and  $\mathcal{C}(\pi)$  are of dimension  $t(\pi) - 1$ . *Proof.* Since subword complexes are pure, then the dimension of the core of the pipe dream complex of  $\pi$  equals the dimension of one of its facets. Consider the facet given by the bottom reduced pipe dream drawn inside the core region. The dimension of this facet equals one less than the number of elbows in the core and from the construction of  $T(\pi)$  this equals  $t(\pi)-2$ . The dimension of the root polytope  $\mathcal{P}(T(\pi))$  is the number of edges in  $T(\pi)$ , which by definition is  $t(\pi)$ .

4.4. BIJECTION BETWEEN REDUCED PIPE DREAMS FOR  $\pi$  AND SIMPLICES OF  $C(\pi)$ . The map M can be easily extended to a map between pipe dreams P of  $\pi$  drawn inside  $R(\pi)$  and forests F on m vertices as follows. For each elbow tile in P add the edge (i, j) corresponding to the box of the elbow to F. Moreover, add the edge (1, m) to F.

THEOREM 4.8. The reduced pipe dreams of  $\pi = 1\pi'$ , where the diagram of  $\pi'$  is a partition, are in bijection with the noncrossing alternating spanning trees of the directed transitive closure of  $T(\pi)$  via the map M.

We prove Theorem 4.8 by induction on the number of columns in the diagram. We break it down in several lemmas.

LEMMA 4.9. Take the permutation  $\pi = 1\pi'$ , where the diagram of  $\pi'$  is  $\lambda = (k)$ . The reduced pipe dreams of  $\pi$  are in bijection with the noncrossing alternating spanning trees of the directed transitive closure of  $T(\pi)$  via the map M.

*Proof.* The edges of  $T(\pi)$  for such a  $\pi$  are (1,2) and (2,j) for  $j = 3, \ldots, k+2$  and thus for the transitive closure of  $T(\pi)$  we add the edges (1,j) with  $j = 3, \ldots, k+2$ . See Figure 12.



FIGURE 12. The region  $R(\pi)$  and the tree  $T(\pi)$  for  $\lambda = (4)$ 

For l = 2, ..., k + 2 let  $T_l$  be a tree on the vertex set [k + 2] consisting of the edges (2, i) for  $2 < i \leq l$  and (1, j) for  $j \geq l$ . Then  $T_l$ , l = 2, ..., k + 2, are all of the noncrossing alternating spanning trees of the directed transitive closure of  $T(\pi)$ . The map M applied to the bottom reduced pipe dream of  $\pi$  yields  $T_{k+2}$ . Moreover, after performing  $0 \leq i \leq k$  ladder moves (there is only one way to do this) on the bottom reduced pipe dream of  $\pi$ , we obtain a reduced pipe dream whose image under M is  $T_{k+2-i}$ . By Theorem 2.7 these are indeed all of the reduced pipe dreams of  $\pi$  concluding the proof.

LEMMA 4.10. Given  $\pi = 1\pi'$ , where the diagram of  $\pi'$  is a partition and that has more than one nonzero column, let its rightmost (shortest) column be of size k. Then in a reduced pipe dream of  $\pi$  the only configurations of crosses and elbows that can occur in the rightmost column of  $cr(\pi)$  are, as read from above, l crosses and k - l elbows, for  $l = 0, \ldots, k$ .

*Proof.* This follows immediately from Theorem 2.7.

LEMMA 4.11. Let  $\pi = 1\pi'$ , where the diagram of  $\pi'$  is a partition  $\lambda = (\lambda_1, \ldots, \lambda_z)$ that has more than one nonzero column. Consider all reduced pipe dreams of  $\pi$  where the configuration of crosses and elbows in the rightmost column of  $\operatorname{cr}(\pi)$  is set to have l crosses and k - l elbows for a fixed  $0 \leq l \leq k$ . These are in bijection with reduced pipe dreams of the permutation  $1w_l$ , where  $w_l$  has diagram  $(\lambda_1 - (k - l), \lambda_2 - (k - l), \ldots, \lambda_{z-1} - (k - l))$ .

*Proof.* Since the bottom k - l boxes of the rightmost column of  $cr(\pi)$  are elbows, it can be seen using Theorem 2.7 that the k - l rows containing crosses one step to the south and one step to the west of these k - l boxes can never move anywhere.

Moreover, the fixed rows of crosses do not affect the ladder moves we can make on the remaining crosses. This allows us to get exactly the reduced pipe dreams for the permutation 1w, where the diagram of w is the diagram of  $\pi$  after ignoring the fixed rows and shortest column, i.e., w has diagram  $(\lambda_1 - (k - l), \lambda_2 - (k - l), \dots, \lambda_{z-1} - (k - l))$ .

The following example illustrates the lemma above.

EXAMPLE 4.12. Let  $1\pi = [164235]$  and suppose l = 1, i.e., we are fixing one cross in position (1,3) and elbow in position (2,3), see Figure 13. The elbow in entry (2,3) causes row 3 to consist of only crosses. Therefore the reduced pipe dreams for [164235] with a cross in entry (1,3) and elbow in (2,3) correspond with the reduced pipe dreams for [15234]. The diagram of this latter permutation has fewer columns.



(A) Core of the permutation with diagram  $\lambda = (4, 2)$ 

(B) Core of the permutation with diagram  $\lambda = (3)$ 

FIGURE 13. New core after applying the reduction of Lemma 4.11 for l = 1 to the core on the left.

LEMMA 4.13. Given  $\pi = 1\pi'$ ,  $\pi'$  dominant, where the length of the shortest column of the diagram of  $\pi'$  is k, the set S of noncrossing alternating spanning trees of the directed transitive closure of  $T(\pi)$  is a disjoint union  $S = S_0 \sqcup \cdots \sqcup S_k$ , where

(4.1) 
$$S_l = \{T \in S : (m-k, m-j) \notin E(T) \text{ for } j = 0, \dots, l-1\}$$
  
 $\cup \{T \in S : (m-k, m-j) \in E(T) \text{ for } j = l, \dots, k-1\},\$ 

for  $0 \leq l \leq k$ , where m is the number of vertices of  $T(\pi)$ .

Note that m - k is the label on the bottom of the last column of the core of  $\pi$ . Thus,  $S_l$  consists of the noncrossing alternating spanning trees of the directed transitive closure of  $T(\pi)$  that do not contain the edges corresponding to the top l crosses in last column of the core of  $\pi$  and contain the edges corresponding to the bottom k - l elbows in last column of the core.

Proof of Lemma 4.13. This follows immediately from the definition of  $T(\pi)$ .

LEMMA 4.14. Let  $\pi = 1\pi'$ ,  $\pi'$  dominant of shape  $(\lambda_1, \ldots, \lambda_z)$  and let  $1w_l$  be the permutation where  $w_l$  has diagram  $(\lambda_1 - (k - l), \lambda_2 - (k - l), \ldots, \lambda_{z-1} - (k - l))$ . Use Lemma 4.11 to draw the core region of  $1w_l$  inside the core region of  $\pi$ . Then all the edges corresponding to the entries outside the core region of  $1w_l$  in a tree  $T \in S_l$  are forced by the last column.



FIGURE 14. Edge labeling for the core of  $1w_1$  coming from the core of [164235] on the left.

Proof. If  $\lambda_z < \lambda_{z-1}$ , then the edges inside  $\operatorname{cr}(\pi)$  and outside of  $\operatorname{cr}(1w_l)$  are precisely those in the last column and in the (k-l) rows one step to the south of the k-lboxes fixed to be elbows. The boxes in these rows are crosses and thus we conclude that in this case all the edges corresponding to boxes outside  $\operatorname{cr}(1w_l)$  are indeed fixed after fixing the last column. If  $\lambda_z = \lambda_{z-1} = \cdots = \lambda_{z-j} < \lambda_{z-j-1}$ , then aside from the boxes outside of  $\operatorname{cr}(1w_l)$  described in the previous sentence, we also have the jboxes to the left of the top most elbow on the last column. We will show that a tree  $T \in S_l$  must contain the edge corresponding to these boxes. Let  $T \in S_l$ , u be the E step below the leftmost of these boxes and v be the N step to the right of these boxes. Since T is an alternating spanning tree, then v must be adjacent either to uor to a vertex before u. Similarly, u must be adjacent either to v or to a vertex after v. The only way noncrossing is preserved is if (u, v) is an edge of T. We continue in this fashion by looking at the second leftmost box and taking the E step below it and again prove that the edge corresponding to that box is in T.

LEMMA 4.15. The set  $S_l$ ,  $0 \leq l \leq k$ , as in Lemma 4.13 is in bijection with reduced pipe dreams of the permutation  $1w_l$ , where  $w_l$  has diagram  $(\lambda_1 - (k - l), \lambda_2 - (k - l), \dots, \lambda_{z-1} - (k - l))$ , via the map M.

We prove this lemma and Theorem 4.8 together by using induction on the number of columns of the diagram of  $\pi'$ .

Proof of Theorem 4.8 and Lemma 4.15. We use induction on the number of columns of the diagram of  $\pi'$ . The base case for Theorem 4.8 where this diagram contains one column is proved in Lemma 4.9. Notice that in the proof of this lemma the base case for Lemma 4.15 is also proven.

By Lemma 4.13 the noncrossing alternating spanning trees of the directed transitive closure of  $T(\pi)$  can be broken down into the sets  $S_l$ ,  $l = 0, 1, \ldots, k$ . Consider the permutation  $1w_l$  where  $w_l$  has diagram  $(\lambda_1 - (k-l), \lambda_2 - (k-l), \ldots, \lambda_{z-1} - (k-l))$ . By inductive hypothesis, we know that  $1w_l$  satisfies Theorem 4.8, i.e., its reduced pipe dreams are in bijection with the noncrossing alternating spanning trees of the directed transitive closure of  $T(1w_l)$  via the map M. By Lemma 4.14 we have that the noncrossing alternating spanning trees of the directed transitive closure of  $T(1w_l)$ yield the set  $S_l$ . Finally, by Lemma 4.11 we know that the reduced pipe dreams of the permutations  $1w_l$ , as l = 0, 1, ..., k, are in bijection with the reduced pipe dreams of  $\pi$ , concluding the proof.

### 5. Reduced forms in the subdivision algebra and Grothendieck polynomials

In this section we show that Grothendieck polynomials of permutations  $\pi = 1\pi'$ ,  $\pi'$  dominant, are special cases of reduced forms in the subdivision algebra of root polytopes. To this end we start by defining the notions appearing in the previous sentence.

The subdivision algebra of root polytopes  $S(\beta)$  is a commutative algebra generated by the variables  $x_{ij}$ ,  $1 \leq i < j \leq n$ , over  $\mathbb{Q}[\beta]$ , subject to the relations  $x_{ij}x_{jk} = x_{ik}(x_{ij} + x_{jk} + \beta)$ , for  $1 \leq i < j < k \leq n$ . This algebra is called the subdivision algebra, because its relations can be seen geometrically as subdividing root polytopes via Lemma 3.1. The subdivision algebra has been used extensively for subdividing root (and flow) polytopes in [12, 13, 14, 15, 16, 17, 18].

A reduced form of the monomial in the algebra  $S(\beta)$  is a polynomial obtained by successively substituting  $x_{ik}(x_{ij} + x_{jk} + \beta)$  in place of an occurrence of  $x_{ij}x_{jk}$  for some i < j < k until no further reduction is possible. Note that the reduced forms are not necessarily unique.

A possible sequence of reductions in algebra  $S(\beta)$  yielding a reduced form of  $x_{12}x_{23}x_{34}$  is given by

$$x_{12}x_{23}x_{34} \rightarrow x_{12}x_{24}x_{23} + x_{12}x_{34}x_{24} + \beta x_{12}x_{24} \rightarrow x_{24}x_{13}x_{12} + x_{24}x_{23}x_{13} + \beta x_{24}x_{13} + x_{34}x_{14}x_{12} + x_{34}x_{24}x_{14} + \beta x_{34}x_{14} + \beta x_{14}x_{12} + \beta x_{24}x_{14} + \beta^2 x_{14} \rightarrow x_{13}x_{14}x_{12} + x_{13}x_{24}x_{14} + \beta x_{13}x_{14} + x_{24}x_{23}x_{13} + \beta x_{24}x_{13} + x_{34}x_{14}x_{12} + x_{34}x_{24}x_{14} + \beta x_{34}x_{14} + \beta x_{14}x_{12} + \beta x_{24}x_{14} + \beta^2 x_{14}$$
(5.1)

where the pair of variables on which the reductions are performed is in **boldface**. The reductions are performed on each monomial separately.

Given a noncrossing tree T on the vertex set [n], let  $m[T] := \prod_{(i,j)\in T} x_{ij}$ . The canonical reduced form  $\operatorname{Crf}_T(x_{ij} \mid 1 \leq i < j \leq n)$  of m[T] is the reduced form obtained by performing reductions on the tree T from front to back (or back to front) on the topmost edges always. This can of course be translated into an algebraic context as follows. For  $x_{ij} = t_i$  we denote by  $\operatorname{Crf}_T(t_1, \ldots, t_{n-1}) = \operatorname{Crf}_T(x_{ij} \mid 1 \leq i < j \leq n)$ . While the reduced form of a monomial in the subdivision algebra is not necessarily unique, once we set  $x_{ij} = t_i$  it becomes unique. This is the statement of the next theorem which we prove in Appendix A.

THEOREM 5.1. Given a noncrossing tree T on the vertex set [n], let  $R_T(x_{ij} | 1 \leq i < j \leq n)$  be an arbitrary reduced form of m[T]. Let  $R_T(t_1, \ldots, t_{n-1})$  be the reduced form  $R_T(x_{ij} | 1 \leq i < j \leq n)$  when we let  $x_{ij} = t_i$ . Then,

(5.2) 
$$R_T(t_1, \dots, t_{n-1}) = \operatorname{Crf}_T(t_1, \dots, t_{n-1}).$$

We will use the notation  $R_T(\mathbf{t})$  when instead of setting  $x_{ij} = t_i$ , we do the following. Let  $i_1 < \cdots < i_v$  be the vertices of T that have outgoing edges. Therefore, the only  $x_{ij}$ 's appearing in a reduced form must have  $i \in \{i_1, \ldots, i_v\}$ . The reduced form  $\tilde{R}_T(\mathbf{t})$  is then obtained from  $R_T(x_{ij} \mid 1 \leq i < j \leq n)$  by setting  $x_{i_k,j} = t_k$  for  $k \in [v]$  and all  $j \in [n]$ . The following theorem provides a combinatorial way of thinking about double Grothendieck polynomials.

THEOREM 5.2 ([6, 10]). The double Grothendieck polynomial  $\mathfrak{G}_w(\mathbf{x}, \mathbf{y})$  for  $w \in S_n$ , where  $\mathbf{x} = (x_1, \ldots, x_{n-1})$  and  $\mathbf{y} = (y_1, \ldots, y_{n-1})$  can be written as

(5.3) 
$$\mathfrak{G}_w(\mathbf{x}, \mathbf{y}) = \sum_{P \in \operatorname{Pipes}(w)} (-1)^{\operatorname{codim}_{PD(w)}F(P)} w t_{x,y}(P),$$

where  $\operatorname{Pipes}(w)$  is the set of all pipe dreams of w (both reduced and nonreduced), F(P) is the interior face in PD(w) labeled by the pipe dream P,  $\operatorname{codim}_{PD(w)}F(P)$  denotes the codimension of F(P) in PD(w) and  $wt_{x,y}(P) = \prod_{(i,j)\in\operatorname{cross}(P)}(x_i - y_j + x_iy_j)$ , with  $\operatorname{cross}(P)$  being the set of positions where P has a cross.

Note that in the product  $\prod_{(i,j)\in cross(P)}(x_i - y_j + x_iy_j)$  appearing in the statement of Theorem 5.2 we are assuming a certain labeling of rows and columns. Conventionally, rows are labeled increasingly from top to bottom and columns are labeled increasingly from left to right. Also recall that the lowest degree terms of  $\mathfrak{G}_w(\mathbf{x}, \mathbf{y})$ give the Schubert polynomial  $\mathfrak{S}_w(\mathbf{x}, \mathbf{y})$ . Except in Theorem 5.4, we will be working with the single Grothendieck polynomial

$$\mathfrak{G}_w(\mathbf{x}) := \mathfrak{G}_w(\mathbf{x}, \mathbf{0}).$$

In other words, for single Grothendieck polynomials we use the weight  $wt_x(P) = \prod_{(i,j)\in cross(P)} x_i$  instead of  $wt_{x,y}(P) = \prod_{(i,j)\in cross(P)} (x_i - y_j + x_i y_j)$  in equation (5.3). In the spirit of Theorem 5.2, we use the following definition for the  $\beta$ -Grothendieck

polynomial:

(5.4) 
$$\mathfrak{G}_{w}^{\beta}(\mathbf{x}) := \sum_{P \in \operatorname{Pipes}(w)} \beta^{\operatorname{codim}_{PD(w)}F(P)} wt_{x}(P).$$

If we set  $\beta = 0$  in (5.4), then we recover the single Schubert polynomial  $\mathfrak{S}_w(\mathbf{x})$ . Note that if in (5.4) we assume that  $\beta$  has degree -1, while all other variables are of degree 1, then the powers of  $\beta$ 's simply make the polynomial  $\mathfrak{G}_w^\beta(\mathbf{x})$  homogeneous. We chose this definition of  $\beta$ -Grothendieck polynomials, as it will be the most convenient notationwise for our purposes.

THEOREM 5.3. Given  $\pi = 1\pi'$ ,  $\pi'$  dominant, we have that for any reduced form of  $m[T(\pi)]$ 

(5.5) 
$$\tilde{R}_{T(\pi)}(\mathbf{t}) = \left(\prod_{i=1}^{n-1} t_i^{g_i}\right) \mathfrak{G}_{\pi^{-1}}^{\beta}(t_1^{-1}, \dots, t_{n-1}^{-1}),$$

where  $g_i$  is the number of boxes in the *i*th column from the left in  $R(\pi)$ .

A special case of Theorem 5.3 for  $\pi = 1n(n-1)...2$  appears in [9].

Note that for  $\pi = 1\pi'$ ,  $\pi'$  dominant then  $\pi^{-1}$  is also of this form.

We now relate the canonical reduced form to the double Grothendieck polynomial. For  $\beta = -1$  denote the canonical reduced form by  $\operatorname{Crf}_T^{\beta=-1}(x_{ij} \mid 1 \leq i < j \leq n)$ . Before we can state and prove Theorem 5.4 we need to define a map  $\phi$  from the labels (i, j) that the boxes in the region  $R(\pi)$  inherit from the labeling of its boundary (as described in Figure 9) to the conventional labeling where rows are labeled increasingly from top to bottom and columns are labeled increasingly from left to right. We call the former labeling the tree labeling and when unclear which labeling we are talking of we put a T index on it:  $(i, j)_T$ . The map  $\phi$  simply takes the tree label (i, j) to the conventional label  $(\phi_{ij}(i), \phi_{ij}(j))$  of the corresponding box. In the example of Figure 9 we have that  $\phi((1, 6)) = (1, 1), \phi((2, 3)) = (3, 2), \phi((5, 6)) = (1, 4), \phi((4, 5)) = (2, 3),$  and so forth.

THEOREM 5.4. Given  $\pi = 1\pi'$ ,  $\pi'$  dominant, we have that

(5.6) 
$$\operatorname{Crf}_{T(\pi)}^{\beta=-1}\left(x_{ij} = \frac{1}{x_{\phi_{ij}(i)} - y_{\phi_{ij}(j)} + x_{\phi_{ij}(i)}y_{\phi_{ij}(j)}} \mid 1 \leq i < j \leq n\right)$$
  
$$= \left(\prod_{(i,j)_T \in R(\pi)} \frac{1}{x_{\phi_{ij}(i)} - y_{\phi_{ij}(j)} + x_{\phi_{ij}(i)}y_{\phi_{ij}(j)}}\right) \mathfrak{G}_{\pi}(\mathbf{x}, \mathbf{y}).$$

*Proof.* By definition we have that

(5.7) 
$$\operatorname{Crf}_{T(\pi)}(x_{ij} \mid 1 \leq i < j \leq n) = \sum_{G \in \mathcal{L}(T(\pi))} \beta^{\operatorname{codim}_{\mathcal{P}(T(\pi))}\mathcal{P}(G)} wt(G),$$

where  $wt(G) = \prod_{(i,j)\in G} x_{ij}$ ,  $\mathcal{L}(T(\pi))$  denotes the set of graphs corresponding to the terms of the reduced form of  $m[T(\pi)]$ , and  $\mathcal{P}(G)$  denotes the simplex in the canonical triangulation of  $\mathcal{P}(T(\pi))$  corresponding to G. Together with Theorem 4.4 using the map M and Theorem 5.2, we obtain (5.6).

### 6. Volumes and Ehrhart series of root polytopes

In this section we state the two immediate corollaries regarding volumes and Ehrhart series of root polytopes following from Theorem 1.1. Recall that the normalized volume of a *d*-dimensional polytope  $\mathcal{P}$  is *d*! times its usual volume, which is always integral for lattice polytopes.

Papers [12, 13] explore a systematic way of calculating volumes of root polytopes. Indeed, the subdivision algebra and the (canonical) reduced form as defined in the previous section is the main tool. An immediate corollary of the results in [12] (relying on the Reduction Lemma for Root polytopes, Lemma 3.1) is:

LEMMA 6.1 ([12]). For a forest G the normalized volume of  $\mathcal{P}(G)$  is equal to the evaluation of the canonical reduced form  $\operatorname{Crf}_T$  at  $x_{ij} = 1$  for all i, j and  $\beta = 0$ .

Lemma 6.1 together with Theorem 5.3 readily imply:

THEOREM 6.2. Let  $\pi = 1\pi'$ , where  $\pi'$  is a dominant permutation. Then the normalized volume of  $\mathcal{P}(T(\pi))$  is equal to the number of reduced pipe dreams of  $\pi$ . This can be written as

(6.1) 
$$\operatorname{vol}(\mathcal{P}(T(\pi))) = \mathfrak{G}_{\pi}^{\beta=0}(\mathbf{1}).$$

Recall that for a polytope  $\mathcal{P} \subset \mathbb{R}^N$ , the  $t^{th}$  dilate of  $\mathcal{P}$  is  $t\mathcal{P} = \{(tx_1, \ldots, tx_N) \mid (x_1, \ldots, x_N) \in \mathcal{P}\}$ . The number of lattice points of  $t\mathcal{P}$ , where t is a nonnegative integer and  $\mathcal{P}$  is a convex polytope, is given by the *Ehrhart function*  $i(\mathcal{P}, t)$ . If  $\mathcal{P}$  has integral vertices then  $i(\mathcal{P}, t)$  is a polynomial.

In order to state the Ehrhart series of root polytopes we need the following lemma, which follows from the well-known relationship of f- and h-vectors. We note that we take  $h(\mathcal{C}, x) = \sum_{i=0}^{d} h_i x^i$  to be the h-polynomial of a (d-1)-dimensional simplicial complex  $\mathcal{C}$ .

LEMMA 6.3 ([23]). Let C be a (d-1)-dimensional pure simplicial complex homeomorphic to a ball and  $f_i^{\circ}$  be the number of interior faces of C of dimension i. Then

(6.2) 
$$h(\mathcal{C}, \beta + 1) = \sum_{i=0}^{d-1} f_i^{\circ} \beta^{d-1-i}$$

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Lemma 6.1 straightforwardly generalizes to encode the *h*-polynomial of  $\mathcal{P}(G)$  in terms of the evaluation of the canonical reduced form  $\operatorname{Crf}_T$  at  $x_{ij=1}$  for all i, j; see [12, 17] for more details. Lemma 6.4 below states the specialized result for the polytopes  $\mathcal{P}(T(\pi))$  considered in this paper.

LEMMA 6.4 ([17]). For any permutation  $\pi$  the following holds:

$$\mathfrak{G}_{\pi}^{\beta-1}(\mathbf{1}) = h(PD(\pi),\beta)$$

THEOREM 6.5. Let  $\pi = 1\pi'$ , where  $\pi'$  is a dominant permutation. Then

(6.3) 
$$\mathfrak{G}_{\pi}^{\beta-1}(\mathbf{1}) = \sum_{m \ge 0} \left( i \big( \mathcal{P}(T(\pi)), m \big) \beta^m \big) (1-\beta)^{\dim \big( \mathcal{P}(T(\pi)) \big) + 1} \right).$$

*Proof.* Since canonical triangulation of  $\mathcal{P}(T(\pi))$  is unimodular and it is equivalent as a simplicial complex to  $\mathcal{C}^2(\pi)$ , we have [23]

(6.4) 
$$h(\mathcal{C}^2(\pi),\beta) = \sum_{m \ge 0} \left( i \left( \mathcal{P}(T(\pi)), m \right) \beta^m \right) (1-\beta)^{\dim \left( \mathcal{P}(T(\pi)) \right) + 1}.$$

By Theorem 1.1 and Lemma 6.3 we get that  $h(\mathcal{C}^2(\pi),\beta) = h(PD(\pi),\beta)$ . Together with Lemma 6.4 this concludes the proof.

### Appendix A. Uniqueness of t-reduced forms

The aim of this appendix is to prove Theorem 5.1, which states that when we let  $x_{ij} = t_i$  for all *i*, then the reduced form becomes unique. For clarity we call the reduced forms with the substitution  $x_{ij} = t_i$ , the *t*-reduced forms. Since the first appearance of this paper alternative proofs have been found for a more general case of this uniqueness property in [8, 19].

In order to prove Theorem 5.1 we recall several definitions and results from [12].

A reduction on the edges (i, j), (j, k) of a noncrossing graph G is *noncrossing* if the graphs resulting from the reduction are also noncrossing. Analogously we can define noncrossing reductions on m[G].

THEOREM A.1 ([12]). Let T be a noncrossing tree on the vertex set [n]. Performing noncrossing reductions on m[T], regardless of order, we obtain a unique reduced form  $R_T^{noncross}(x_{ij} \mid 1 \leq i < j \leq n)$  for m[T].

Consider a noncrossing tree T on [n]. We define the *pseudo-components* of T inductively. The unique simple path P from 1 to n is a pseudo-component of T. The graph  $T \setminus P$  is an edge-disjoint union of trees  $T_1, \ldots, T_k$ , such that if v is a vertex of P and  $v \in T_l$ ,  $l \in [k]$ , then v is either the minimal or maximal vertex of  $T_l$ . Furthermore, there are no k - 1 trees whose edge-disjoint union is  $T \setminus P$  and that satisfy all the requirements stated above. The set of pseudo-components of T, denoted by ps(T) is  $ps(T) = \{P\} \cup ps(T_1) \cup \cdots \cup ps(T_k)$ . A pseudo-component P' is said to be on [i, j], i < j if it is a path with endpoints i and j. A pseudo-component P' on [i, j] is said to be a *left pseudo-component* of T if there are no edges  $(j, s) \in E(T)$  with s < i and a right pseudo-component if there are no edges  $(j, s) \in E(T)$  with j < s. See Figure 15 for an example.

THEOREM A.2 ([12]). Let T be a noncrossing tree. Then  $R_T^{noncross}(x_{ij} | 1 \leq i < j \leq n)$ is the sum of the monomials corresponding to the following graphs weighted with powers of  $\beta$  (of degree 1) to obtain a homogeneous polynomial. The graphs are: all noncrossing alternating spanning forests of the directed transitive closure of T on the vertex set [n] containing edge (1, n) and at least one edge of the form  $(i_1, j)$  with  $i_1 \leq i$ for each right pseudo-component of T on [i, j] and at least one edge of the form  $(i, j_1)$ with  $j \leq j_1$  for each left pseudo-component of T on [i, j].



FIGURE 15. The edge sets of the pseudo-components in the graph depicted are  $\{(1,5), (5,8)\}, \{(2,5)\}, \{(3,4), (4,5)\}, \{(5,6), (6,7)\}$ . The pseudo-component with edge set  $\{(1,5), (5,8)\}$  is a both a left and right pseudo-component, while the pseudo-components with edge sets  $\{(2,5)\}, \{(3,4), (4,5)\}$  are left pseudo-components and the pseudo-component with edge set  $\{(5,6), (6,7)\}$  is a right pseudo-component.

We note that in the above we assume that the vertices of our graphs are drawn on a line from left to right in increasing order, 1, 2, ..., n. This condition is of course not an essential condition for the above theorems, and if we rearrange the order of the vertices of our graphs, then we can reinterpret the above results accordingly.

Consider the noncrossing tree T on the vertex set [n] with vertices drawn from left to right in increasing order 1, 2, ..., n. Let (k, l), (l, m) be a pair of nonalternating edges in T. If the reduction performed on (k, l), (l, m) is noncrossing, then we set  $T_{klm} = T$  with  $T_{klm}$  drawn identically to T. If the reduction performed on (k, l), (l, m)is not noncrossing, then let  $C_1 = (V_1, E_1)$  and  $C_2 = (V_2, E_2)$  be the connected components containing the vertices k and m, respectively, in the graph  $T - \{(k, l), (l, m)\} =$  $([n], E(T) \setminus \{(k, l), (l, m)\})$ . Then we define  $T_{klm} = T$  to be drawn with its vertices arranged from left to right in the following order:  $v_1^1, \ldots, v_1^p, w_1, \ldots, w_q, v_2^1, \ldots, v_2^r$ , where  $V_1 = \{v_1^1 < \cdots < v_1^p\}, V_2 = \{v_2^1 < \cdots < v_2^r\}, [n] \setminus (V_1 \cup V_2) = \{w_1 < \cdots < w_q,\}$ . The tree  $T_{klm}$  is then a noncrossing tree. See Figure 16.



FIGURE 16. Reduction performed on the edges (2, 4), (4, 5) of T is not noncrossing, however when performed on  $T_{245}$  it is noncrossing.

LEMMA A.3. For a noncrossing tree T on the vertex set [n] and any two edges (k,l), (l,m) of T that are nonalternating, we have that

(A.1) 
$$R_T^{noncross}(t_i \mid 1 \leq i \leq n-1) = R_{T_{klm}}^{noncross}(t_i \mid 1 \leq i \leq n-1).$$

*Proof.* We prove this lemma by induction on the number of increasing paths in T. Suppose there is a vertex  $v \neq l$  that is nonalternating. Perform noncrossing reductions at v in both T and  $T_{klm}$  obtaining three descendants. Note that the graphs obtained in this fashion from T and  $T_{klm}$  are in natural bijection, and they each have fewer number of increasing paths than does T, so by inductive hypothesis the lemma is true for them. However,  $R_T^{noncross}(t_i \mid 1 \leq i \leq n-1)$  and  $R_{T_{klm}}^{noncross}(t_i \mid 1 \leq i \leq n-1)$ is the sum of the t-reduced forms corresponding to the mentioned graphs, so we are done.

It remains to prove the case when the only nonalternating vertex of T is l. This is accomplished in Lemma A.4 below.

LEMMA A.4. For  $T := T^l = ([n], \{(i,l), (l,j) \mid 1 \leq i < l, l < j \leq n\})$ , for some  $2 \leq l \leq n-1$ , and any two edges (k,l), (l,m) of T that are nonalternating, we have that

(A.2) 
$$R_T^{noncross}(t_i \mid 1 \leq i \leq n-1) = R_{T_{klm}}^{noncross}(t_i \mid 1 \leq i \leq n-1).$$

*Proof.* The only both left and right pseudo-component of T is  $\{(1,l), (l,n)\}$ , its left pseudo-components are  $\{(i,l) \mid 2 \leq i \leq l-1\}$ , and its right pseudo-components are  $\{(l,i) \mid l+1 \leq i \leq n-1\}$ . Similarly, the only both left and right pseudo-component of T is  $\{(k,l), (l,m)\}$ , its left pseudo-components are  $\{(i,l) \mid i \neq k, 1 \leq i \leq l-1\}$ , and its right pseudo-components are  $\{(l,i) \mid i \neq m, l+1 \leq i \leq n\}$ . Using this one can prove by induction on l that there is a bijection between the forests described in Theorem A.2 for T and for those of  $T_{klm}$  such that the number of edges emanating from any vertex  $i \in [n]$  is preserved. While such a proof is not hard, it is technical to describe, and we leave it to the interested reader.

Proof of Theorem 5.1. We proceed by induction on the number of increasing paths in T. If we start by a noncrossing reduction (k, l), (l, m) in T, then no matter how we reduce the three descendants of T which each have fewer number of increasing paths, we obtain that the t-reduced form of m[T] we get is  $\operatorname{Crf}_T(\mathbf{t})$ .

Suppose we start with a reduction (k, l), (l, m) in T that is a crossing reduction. Redraw the tree T as  $T_{klm}$ . Since we can apply the inductive hypothesis to all three descendants of  $T_{klm}$ , we get that the t-reduced form of m[T] obtained this way is  $\operatorname{Crf}_{T_{klm}}(\mathbf{t})$ .

However, by Lemma A.3  $\operatorname{Crf}_T(\mathbf{t}) = \operatorname{Crf}_{T_{klm}}(\mathbf{t})$ , thereby proving the theorem.  $\Box$ 

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